

22-26 October 2017

Introducing an almost reliable UDP protocol: The Keyed UDP

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Does this looks familiar?



What about this?

Daily average IPv4 packet size vs. date



Quick overview of the Internet traffic in the last 10 years



evolution++

Recorded number of IPv6 packets vs. date







Daily average IPv6 packet size vs. date



And what about applications?

- We have witnessed:
 - more social media
 - more Cloud
 - more X as a Service
 - (virtualization of desktop on the User Side Resources to the Cloud - why?)

Applications in the future

- What can we expect to see in the future?
 - mobile phones now host octacores
 - mobile operators sell data packets
 - every device has some Web App or some Internet service connection (*e.g.* problems with hacked cars)
 - my wristwatch monitors me, my car monitors me, my TV monitors me, my refrigerator monitors everyone, ...
 - "every device will be a screen to The One" (Kevin Kelly, The first 5000 days of the Web)

Applications in the future

- User driven (social media, games, apps...)
- Government driven (eGovernment, eCitizenship, ...)
- Corporations driven (eCustomer)
- Machine driven (P2P, machines that are Internet connected)
- Open source hardware (Arduíno, Raspberry Pi-Like, etc.)
- Open source software (Linux, Python, etc.)

What can we expect?



What can we do?

- Problem 1: The Ethernet payload limit (99.999% of the measured packets that travel in the Internet are < 1500B)
- Problem 2: The TCP control mechanism (Three Way Handshake, AIMD)
- Problem 3: low efficiency of transport protocols (high overhead)
- Problem 4: increase in the size of messages
- Problem 5: transition to IPv6
- Problem 6: increased mobile traffic

Really, what can we do?

 We need new protocols for Machine-to-Machine communication, for real time data transmission, which integrate data imputation, resiliency on the software side, not on the network side

New Keyed-UDP

- UDP is faster and lighter than TCP
- Let us use UDP in a manner that sender and/or receiver can know, to a degree of certainty, if there was packet loss and what packets were lost
- I call this Keyed UDP

 In any normal UDP or TCP communication between two apps (or two machines), the data flows from one port-IP address pair to another port-IP address pair.











IP Addr. 1.1.1.1

IP 1.1.1.1 Dst IP 2.2.2.2 Src Port: 10000 Dest Port: 700

2.2.2.2

7004

- sKUDP: from many ports to one port
- dKUDP: from one port to many ports
- sdKUDP: from many ports to many ports

 Ports don't have to be sequential, *i.e.*, the key can be 7000, 5000, 7001, 7000, 5001, 7000

• So how can the key length (and type) be detected?

Key is			
Protocol	Determined	Discovered	Agreed
dKUDP	Hardcoded	n.a.	Key Definition Protocol
sKUDP	Hardcoded	Inference	Key Definition Protocol
sdKUDP	Hardcoded	Inference	Key Definition Protocol

Retro-compatibility with non-KUDP apps?

Protocol Which application is UDP	dKUDP	sKUDP
Source	Destination application receives all	Destination application receives all
application	packets to a single destination port,	packets from a single source port,
is standard	disabling out-of-order and loss	disabling out-of-order and loss event
UDP	event detection	detection
Destination	Only every nth packet will be	This is almost the case for the
application	received because the source	standard UDP scenario, so no loss
is standard	application will try to send packets	nor out-of-order event detection is
UDP	to destination ports that are not	possible
	being monitored by the destination application	

Keyed-IPv6

- What if instead of using port numbers as keys, we use IPv6 addresses?
 - it may be feasible, most operating systems allow multiple IPv6 address assignment to an interface...

- Pending issues:
 - What to do when the receiving app detects losses? (NOP, report back, ask to resend, data imputation - it will depend on the app)
 - Increase the complexity of the loss-switch inference algorithm by timestamping the packets at arrival?
 - Problems with non-homomorphic NAT-PT machines that are not sufficiently persistent?
 - NAT-PT tables overload? / IPv6 routing tables overload?
 - Others?

• Proposal for a reconstruction algorithm



Keyed-UDP



Rec.ved	ltera	ation>	1		2		3		4		5		6		7		8		9		10	
1a		1a	1a	1a		1a		1 a		1a		1 a		1 a		1a		1 a		1 a		1a
3a		3a	f	3a	2b																	
4a		4a	3a	4a	3a	4a	f													4a		4a
6a		6 a	4a	6a	4a	6 a	4a	6a	4b													
5a		5a	5a	5a	5a	5a	5a	5a	5a	5a	5a											
2b		2b	6a	2b	6a	2b	6a	2b	6a	2b	6b	2b	6b			2b				2b		2b
4b						4b	f	4b	f	4b	f	4b	f	4b	f							
6b							2b	6b	2b	6b	f	6b	2b	6b	2c	6b	2c					
3b							f		f	3b	3b	3b	3b	3b	3c	3b	3b	3b	3b			
2c		2c				2c	4b	2c	f		4b	2c	4b	2c	4b	2c	f	2c	4c	2c	4c	
3c				3c					f				f	3c	f	3c	f	3c	f	3c	f	3c
1c									6b				f		6b	1c	6b	1c	f	1c	6c	1c
4c													f		f		1c	4c	1c	4c	1c	4c
6c				6c									2c		f		f		2c	6c	2c	6c
1d		1d		1d		1d		1d		1d		1d			3c	1d	3c	1d	3c	1d	3c	1d
2d		2d				2d		2d				2d		2d						2d		
4d																				4d		
3d																						
5d						5d				5d		5d						5d				

For the 1^{st} packet there is only one candidate, this is packet 1a, therefore, packet 1a takes the first position.

For the 3^{rd} packet the candidates are $\{3a, 3a, f\}$, therefore, packet 3a wins.

For the 2nd position, the candidates are $\{f, 2b\}$. As there is a tie between 2b and f, the first candidate wins, *i.e.*, the algorithm concludes that the 2nd packet never arrived.

For the 4th position, the candidates are $\{4a, 4a, 4a, 4b\}$, therefore, 4a wins.

The candidates for the remaining positions (5 to 10) are, respectively:

5th: {5a, 5a, 5a, 5a, 5a}, elected 5a
6th: {6a, 6a, 6a, 6a, 6b, 6b}, elected 6a
7th: {f, f, f, f}, elected f

8th: {*2b, 2b, f, 2b, 2c, 2c*}, elected *2b*

9th: {*f*, *f*, *3b*, *3b*, *3c*, *3b*, *3b*}, elected *3b*

10th: {*4b, f, 4b, 4b, 4b, f, 4c, 4c*}, elected *4b*.

Rec.ved	Iteration	> 1		2		3		4		5		6		7		8		9		10	
1a	1a	1a																			
3a	За	f	3a	2b																	
4a	4a	3a	4a	3a	4a	f															
6a	6a	4a	6a	4a	6a	4a	6a	4b													
5a	5a	5a	5a	5a	5a	5a	5a	5a	5a	5a											
2b		6a	2b	6a	2b	6a	2b	6a	2b	6b	2b	6b									
4b					4b	f															
6b						2b	6b	2b	6b	f	6b	2b	6b	2c	6b	2c					
3b						f		f	3b	3b	3b	3b	3b	3c	3b	3b	3b	3b			
2c						4b		f		4b	2c	4b	2c	4b	2c	f	2c	4c	2c	4c	
3c								f				f	3c								
1c								6b				f		6b	1c	6b	1c	f	1c	6c	1c
4c												f		f		1c	4c	1c	4c	1c	4c
6c												2c		f		f		2c	6c	2c	6c
1d														3c		3c		3c		3c	1d
2d																					
4d																					
3d																					
5d																					

1	Ts=																
{1a 1d	2a 2d	3a <mark>3d</mark>	4a 4d	<mark>5a</mark> 5d	6a 6d	1b 1e	2b 2e	<mark>3b</mark> 3e	4b 4e	5b 5e	6b 6e	<mark>1c</mark> }	2c	3с	4c	5c	6с
	Rs=																
{1a	3a		4a	6a		5a	2b	4b		6b	3b	20	C	3c	<u>1c</u>		4c
6c	1d		2d	4d		3d	5d	6d		1e	3e	40	е	6e	<mark>5e</mark> }		
	Fs=																
{1a	f	3a	4a	5a	6a	f	2b	3b	4b	f	6b	1c	2c	3c	4c	5d	6c
1d	2d	3d	4d	5d	6d	1e	f	Зе	4e	5e	6e	}.					
1	6.79	%	OSS	+ 2	209	% S\	witc	hes	5								

Keyed-UDP



Finally, Users!

- We will be flooded with data from all of our machines
- We won't be able to process it, nor we will be aware of its existence
- We will see multi-sensor, multi-location, multi-device, multi-format data only by its impact in our daily lives (a bit like targeted advertising on the email systems)
- After the ubiquity of computers, we will have the ubiquity of data, and for users, this will mean that we will need systems that receive, summarise, process, extract relevant features and feed the appropriate algorithms for our own benefit
- The following step is to use the data from multiple users to infer and confirm daily living patterns, as to allow its monitoring and training

Fueling the Internet of Everything

- will be the Internet of our private lives, our machines, our health and our social interactions
- is this scary? a lot!
- is this worth the effort? probably yes!

Questions!



This work was supported by National Funding from the FCT -Fundação para a Ciência e a Tecnologia, through the UID/EEA/50008/2013 Project. The author acknowledges the support of COST Action IC1303 AAPELE - Architectures, Algorithms and Protocols for Enhanced Living Environments

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